



上海交通大学  
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# Chapter 3: Dynamic Programming

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# Reading

Cormen book:

Thomas, H. ,Cormen, Charles, E., Leiserson, and Ronald, L., Rivest .  
Introduction to Algorithms, The MIT Press.

(read Chapter 16 and 17, page 299-355).



# Elements of dynamic programming

- Two elements are required
  1. Optimal substructure
    - An optimal solution contains within it optimal solutions to the subproblems
  2. Overlapping subproblems
    - Recursive formula exists



# Needleman/Wunsch global alignment (1970)

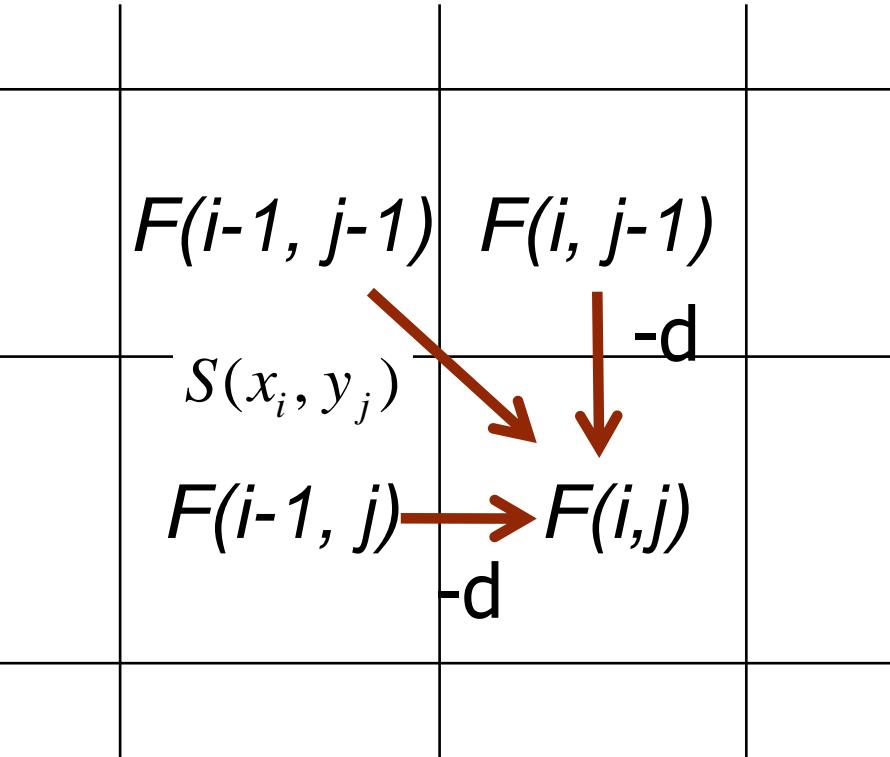
- Two sequences  $X = x_1 \dots x_n$  and  $Y = y_1 \dots y_m$
- Let  $F(i, j)$  be the optimal alignment score of  $X_{1 \dots i}$  of  $X$  up to  $x_i$  and  $Y_{1 \dots j}$  of  $Y$  up to  $y_j$  ( $0 \leq i \leq n, 0 \leq j \leq m$ ), then we have

$$F(0,0) = 0$$

$$F(i, j) = \max \begin{cases} F(i-1, j-1) + s(x_i, y_j) \\ F(i-1, j) - d \\ F(i, j-1) - d \end{cases}$$



# Needleman/Wunsch global alignment (1970)



$$F(0,0) = 0$$

$$F(i, j) = \max \begin{cases} F(i-1, j-1) + s(x_i, y_j) \\ F(i-1, j) - d \\ F(i, j-1) - d \end{cases}$$

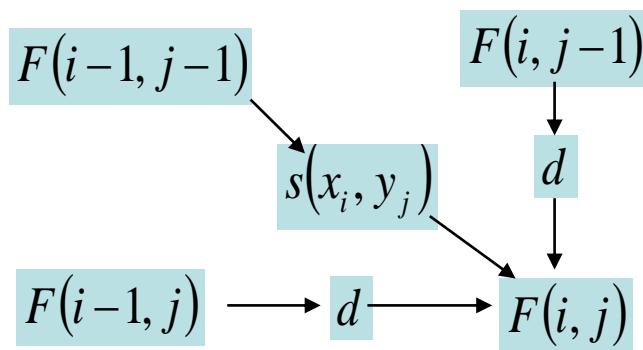


# A simple example

	A	C	G	T
A	2	-7	-5	-7
C	-7	2	-7	-5
G	-5	-7	2	-7
T	-7	-5	-7	2

Find the optimal alignment of AAG and AGC.  
Use a gap penalty of  $d=-5$ .

		A	A	G
A				
G				
C				

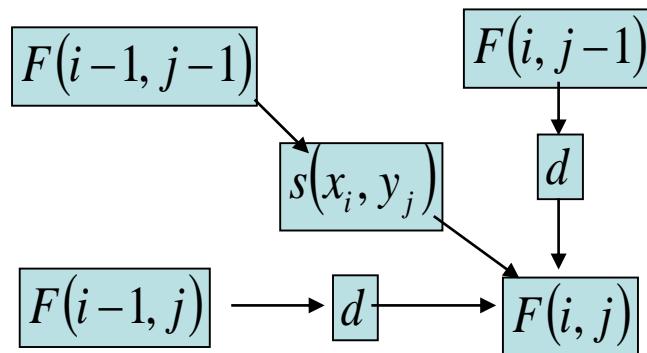




# A simple example

	A	C	G	T
A	2	-7	-5	-7
C	-7	2	-7	-5
G	-5	-7	2	-7
T	-7	-5	-7	2

Find the optimal alignment of AAG and AC  
Use a gap penalty of  $d=-5$ .

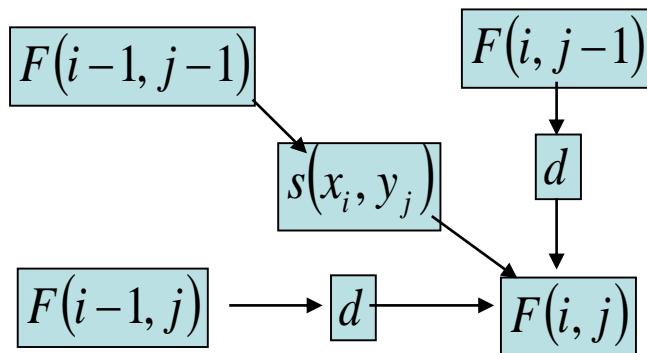


		A	A	G
	0			
A				
G				
C				



# A simple example

	A	C	G	T
A	2	-7	-5	-7
C	-7	2	-7	-5
G	-5	-7	2	-7
T	-7	-5	-7	2



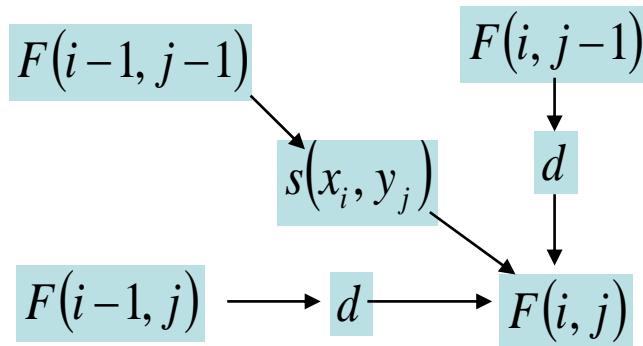
Find the optimal alignment of AAG and ACCTG  
Use a gap penalty of  $d=-5$ .

		A	A	G
	0 →	-5 →	-10 →	-15
A	-5			
G	-10			
C	-15			



# A simple example

	A	C	G	T
A	2	-7	-5	-7
C	-7	2	-7	-5
G	-5	-7	2	-7
T	-7	-5	-7	2



Find the optimal alignment of AAG and ACCTG  
Use a gap penalty of  $d=-5$ .

		A	A	G
	0 →	-5 →	-10 →	-15
A	↓ -5	2 →	-3 →	-8
G	↓ -10	↓ -3	-3 →	-1
C	↓ -15	↓ -8	↓ -8	-6



# Traceback

- Start from the lower right corner and trace back to the upper left.
- Each arrow introduces one character at the end of each aligned sequence.
- A horizontal move puts a gap in the left sequence.
- A vertical move puts a gap in the top sequence.
- A diagonal move uses one character from each sequence.



# A simple example

- Start from the lower right corner and trace back to the upper left.
- Each arrow introduces one character at the end of each aligned sequence.
- A horizontal move puts a gap in the left sequence.
- A vertical move puts a gap in the top sequence.
- A diagonal move uses one character from each sequence.

Find the optimal alignment of AAG and AGC  
Use a gap penalty of  $d=-5$ .

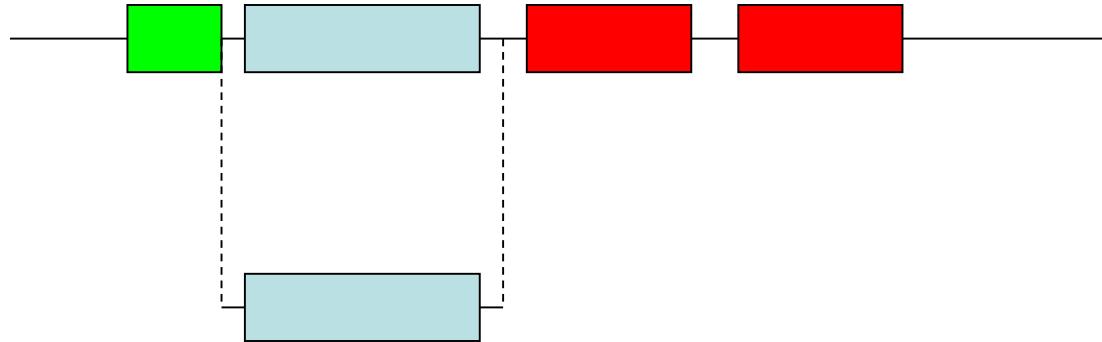
		A	A	G
	0	-5		
A		2	-3	
G				-1
C				-6

AAG-  
-AGC

AAG-  
A-GC



# Local alignment



- Ⓐ A single-domain protein may be homologous to a region within a multi-domain protein.
- Ⓑ Usually, an alignment that spans the complete length of both sequences is not required.



# Smith/Waterman local alignment (1981)

- Two sequences  $X = x_1 \dots x_n$  and  $Y = y_1 \dots y_m$
- Let  $F(i, j)$  be the optimal alignment score of  $X_{1 \dots i}$  of  $X$  up to  $x_i$  and  $Y_{1 \dots j}$  of  $Y$  up to  $y_j$  ( $0 \leq i \leq n, 0 \leq j \leq m$ ), *then we have*

$$F(0,0) = 0$$

$$F(i, j) = \max \begin{cases} 0 \\ F(i-1, j-1) + s(x_i, y_j) \\ F(i-1, j) - d \\ F(i, j-1) - d \end{cases}$$



# Local alignment

- Two differences with respect to global alignment:
  - No score is negative.
  - Traceback begins at the highest score in the matrix and continues until you reach 0.
- Global alignment algorithm: *Needleman-Wunsch*.
- Local alignment algorithm: *Smith-Waterman*.

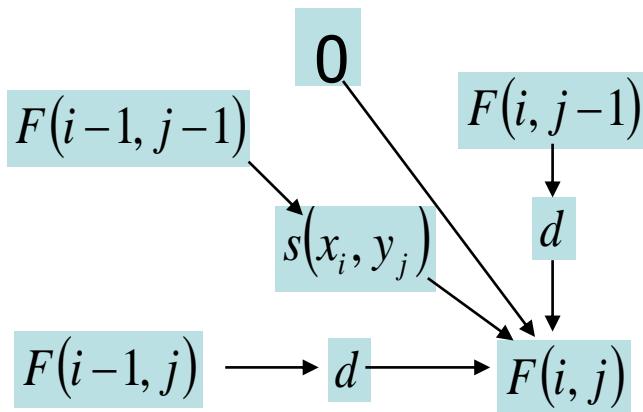


# A simple example

	A	C	G	T
A	2	-7	-5	-7
C	-7	2	-7	-5
G	-5	-7	2	-7
T	-7	-5	-7	2

Find the optimal local alignment of AAG and AGC.  
Use a gap penalty of  $d=-5$ .

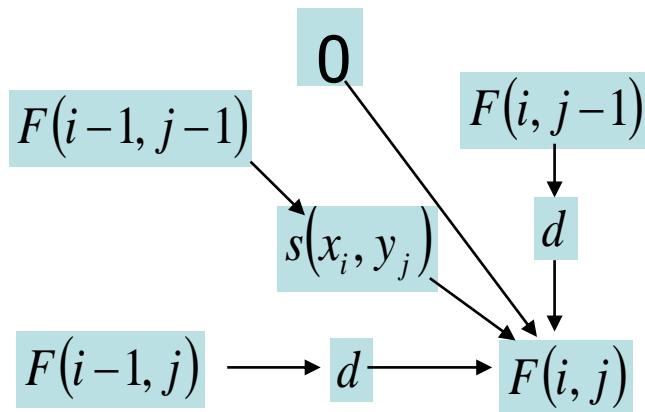
		A	A	G
A				
G				
C				





# A simple example

	A	C	G	T
A	2	-7	-5	-7
C	-7	2	-7	-5
G	-5	-7	2	-7
T	-7	-5	-7	2



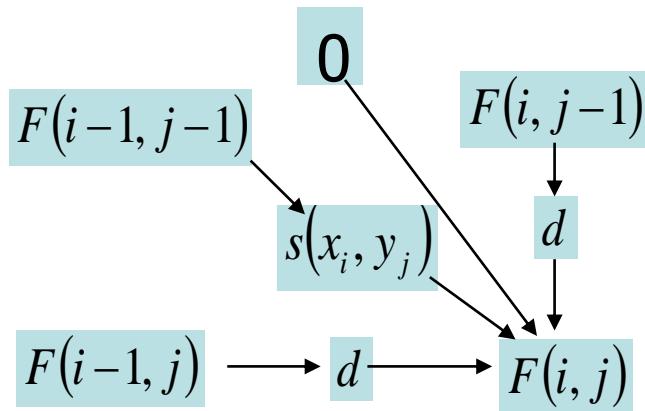
Find the optimal local alignment of AAG and AGC.  
Use a gap penalty of  $d=-5$ .

		A	A	G
	0	0	0	0
A	0			
G	0			
C	0			



# A simple example

	A	C	G	T
A	2	-7	-5	-7
C	-7	2	-7	-5
G	-5	-7	2	-7
T	-7	-5	-7	2



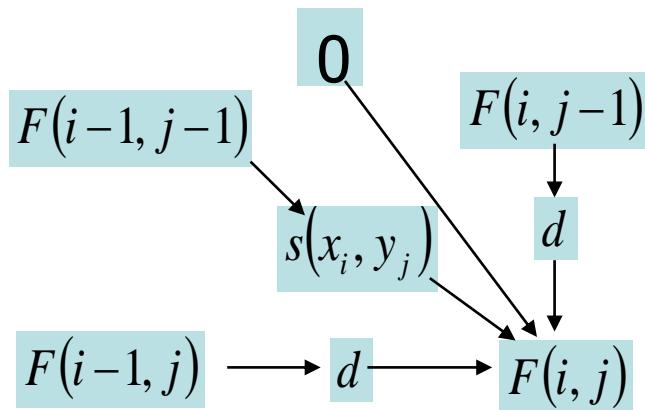
Find the optimal local alignment of AAG and AGC.  
Use a gap penalty of  $d=-5$ .

		A	A	G
	0	0	0	0
A	0	2	2	0
G	0	0	0	4
C	0	0	0	0



# A simple example

	A	C	G	T
A	2	-7	-5	-7
C	-7	2	-7	-5
G	-5	-7	2	-7
T	-7	-5	-7	2



Find the optimal local alignment of AAG and AGC.  
Use a gap penalty of  $d=-5$ .

		A	A	G
	0	0	0	0
A	0	2	2	0
G	0	0	0	4
C	0	0	0	0

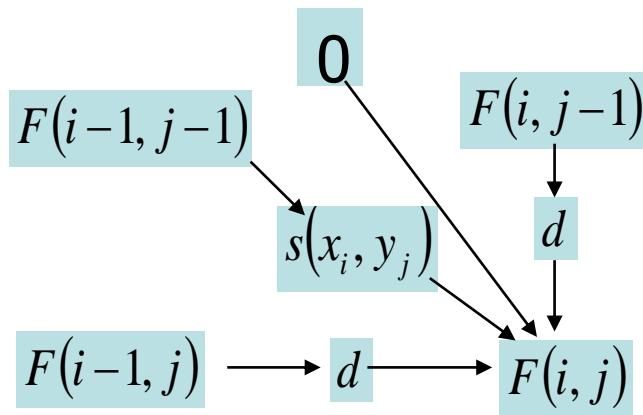
AG

AG



# Local alignment

	A	C	G	T
A	2	-7	-5	-7
C	-7	2	-7	-5
G	-5	-7	2	-7
T	-7	-5	-7	2



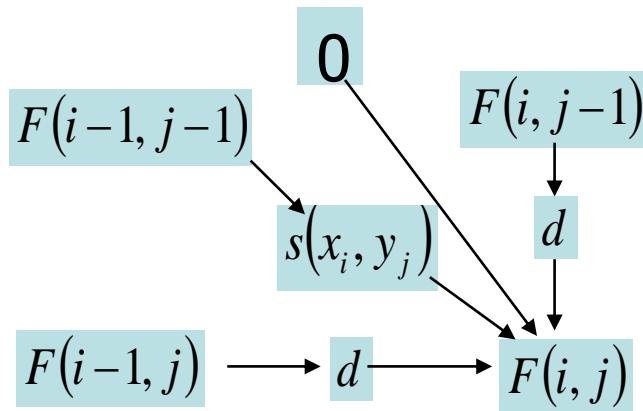
Find the optimal local alignment of AAG and GAAGGC.  
Use a gap penalty of  $d=-5$ .

		A	A	G
	0	0	0	0
G	0			
A	0			
A	0			
G	0			
G	0			
C	0			



# Local alignment

	A	C	G	T
A	2	-7	-5	-7
C	-7	2	-7	-5
G	-5	-7	2	-7
T	-7	-5	-7	2



Find the optimal local alignment of AAG and GAAGGC.  
Use a gap penalty of  $d=-5$ .

		A	A	G
	0	0	0	0
G	0	0	0	2
A	0	2	2	0
A	0	2	4	0
G	0	0	0	6
G	0	0	0	2
C	0	0	0	0



# Greedy algorithm: Choose the best at the moment

- Not always produce the optimal result
- Two elements are required to find an optimal solution by greedy algorithm
  1. Greedy-choice property
    - Global optimal can be reached by local optimal (greedy)
  2. Optimal substructure
    - An optimal solution contains within it optimal solutions to the subproblems



# Acknowledgement

PPTs for examples in dynamic programming are kindly provided by Dr. Qi Liu.